

# FIFO WITH OFFSETS

***HIGH SCHEDULABILITY WITH LOW OVERHEADS***

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# FIFO SCHEDULING

First-In-First-Out (**FIFO**)  
scheduling

extremely simple

very low overheads

*ideal for:*

*IoT-class devices*

*deeply embedded systems*

*hardware implementations*



very **low schedulability**

meeting **deadlines?**



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## THIS PAPER

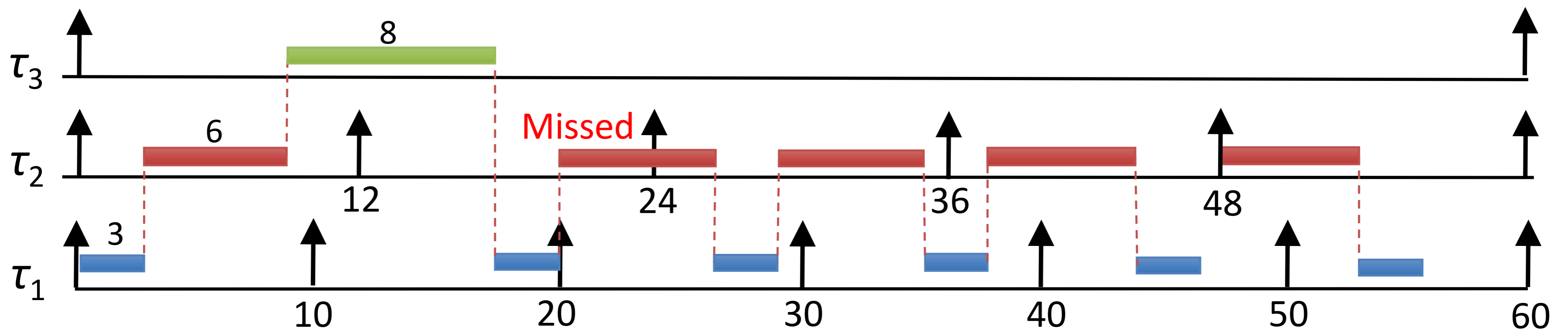
***FIFO can actually achieve excellent schedulability!***

[periodic non-preemptive tasks on a uniprocessor]

**INTUITION**

# THE PROBLEM WITH **PLAIN FIFO** SCHEDULING

**FIFO** schedule of 3 periodic tasks:

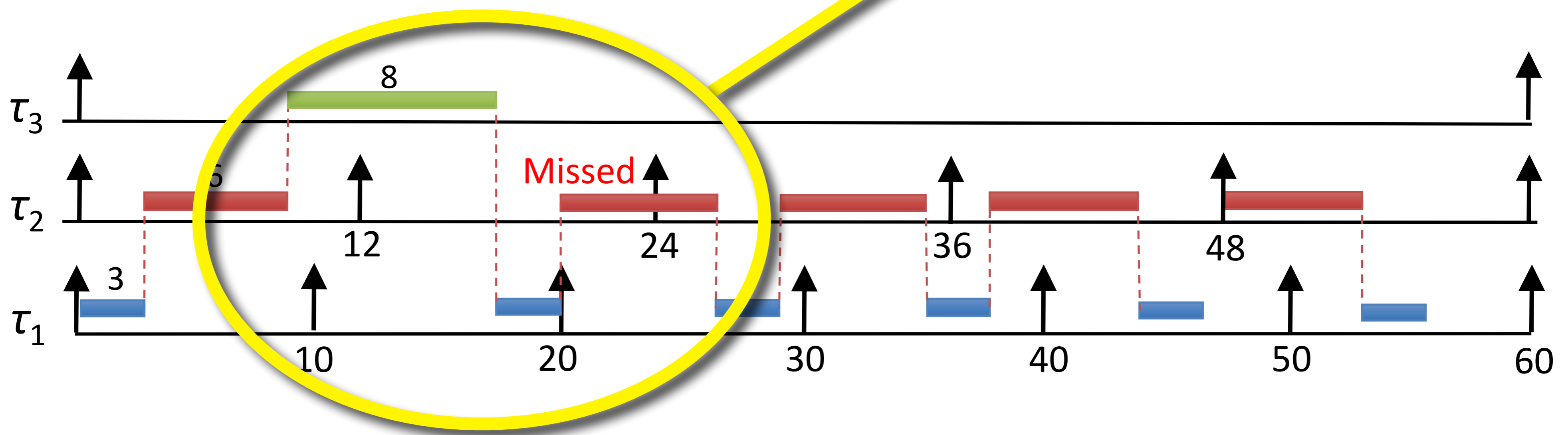


Task	WCET	Period
$\tau_3$	8	60
$\tau_2$	6	12
$\tau_1$	3	10

# THE PROBLEM

**Plain FIFO** is oblivious to deadlines and priorities  
 $\tau_3$  comes first  $\rightarrow$  *deadline miss*

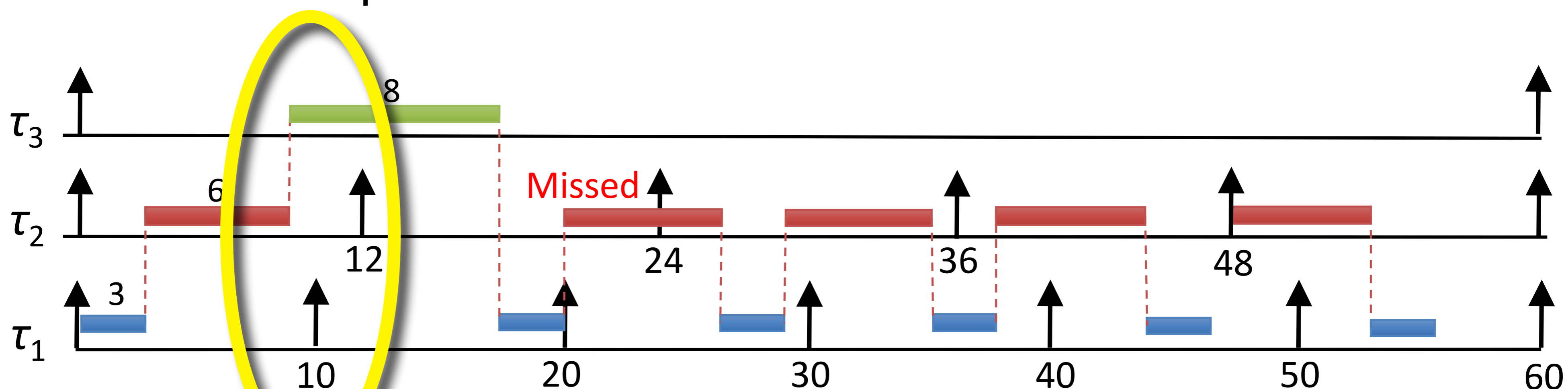
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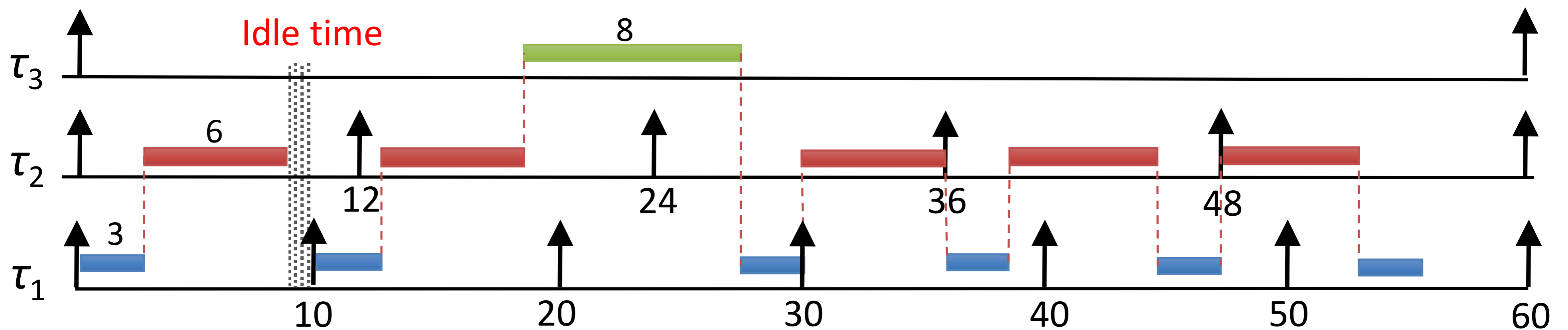


**In fact, *any work-conserving policy* (EDF, RM, ...) must schedule  $\tau_3$  here  $\rightarrow$  deadline miss.**

# NON-WORK-CONSERVING SCHEDULING

[critical-window EDF: Nasri & Fohler, 2016]

**CW-EDF** schedule of the same 3 periodic tasks:



Task	WCET	Period
$\tau_3$	8	60
$\tau_2$	6	12
$\tau_1$	3	10

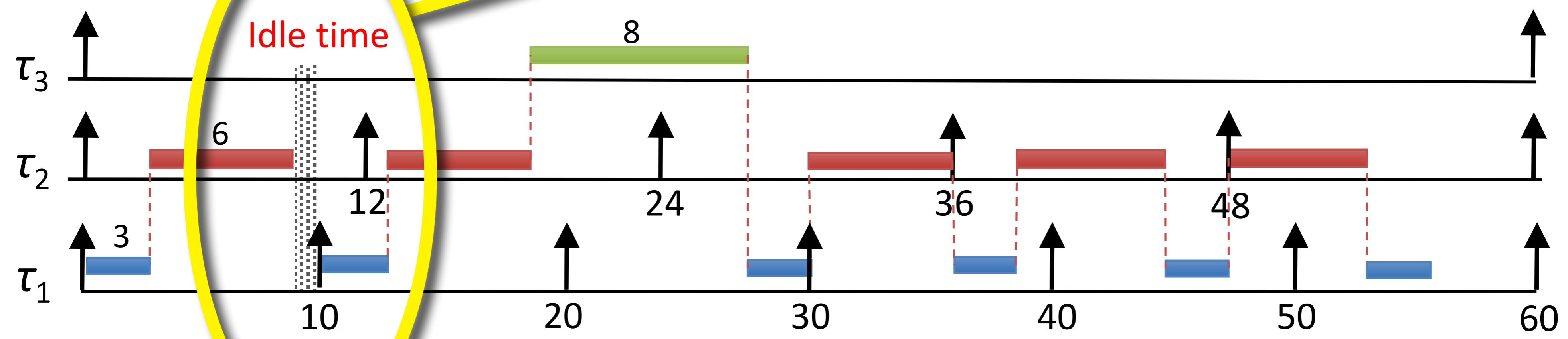


# NON-WORKING

**CW-EDF considers *future job arrivals* in the "critical window" and **postpones**  $\tau_3$  until later.**

[critical-window EDF: Nasri & ...]

**CW-EDF** schedule of the same 3 periodic tasks.

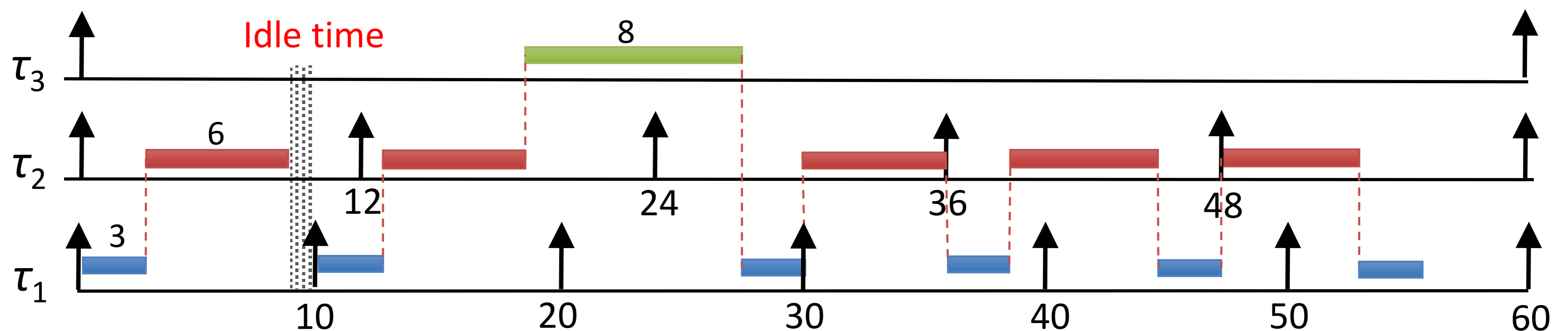


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# NON-WORK-CONSERVING SCHEDULING

[critical-window EDF: Nasri & Fohler, 2016]

**CW-EDF** schedule of the same 3 periodic tasks:



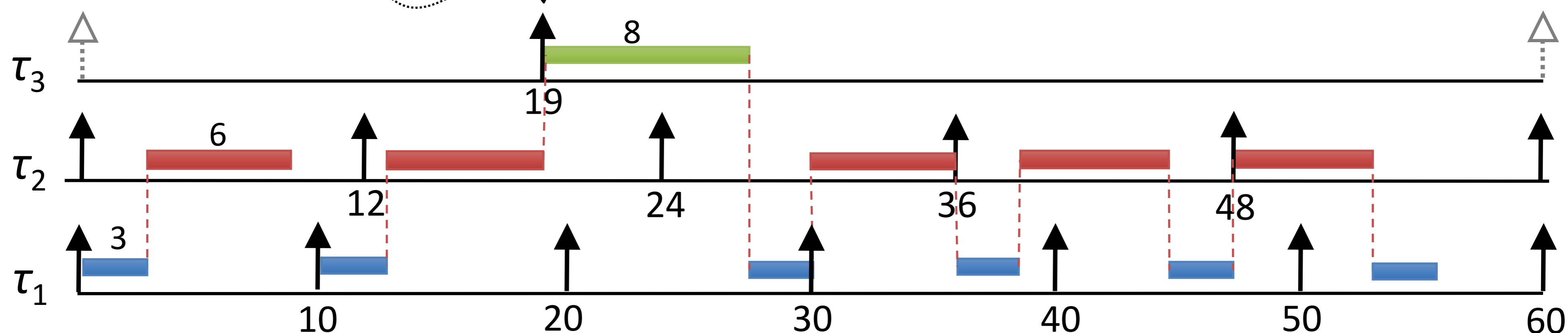
## LIMITATION

CW-EDF incurs *much higher runtime overheads* than simple work-conserving policies.

*ATMega2560 @ 16 MHz: 9.2× higher than RM!*

# INTUITION: FIFO + "JUST THE RIGHT" OFFSETS

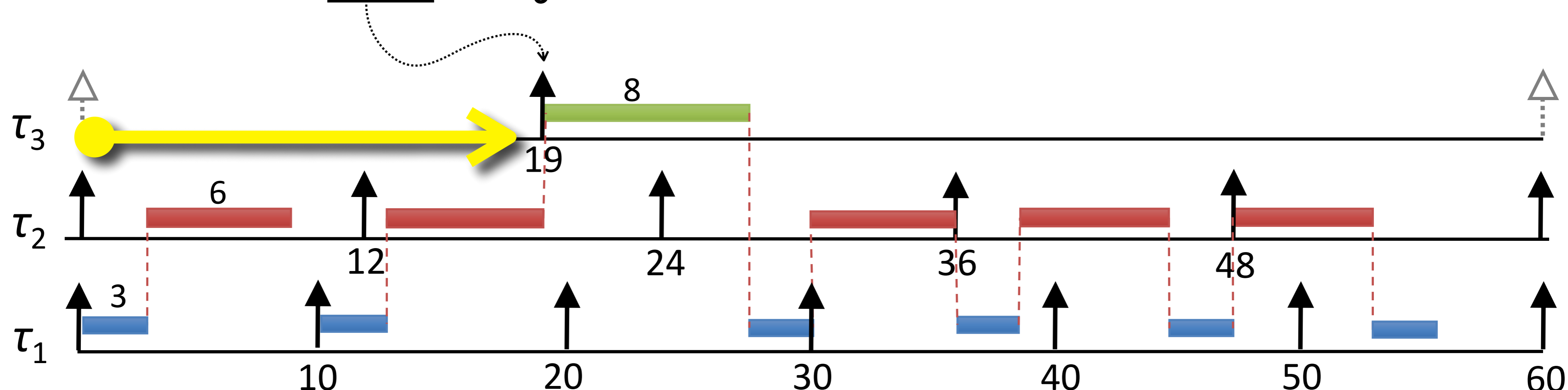
**FIFO** schedule + *offset* for  $\tau_3$ :



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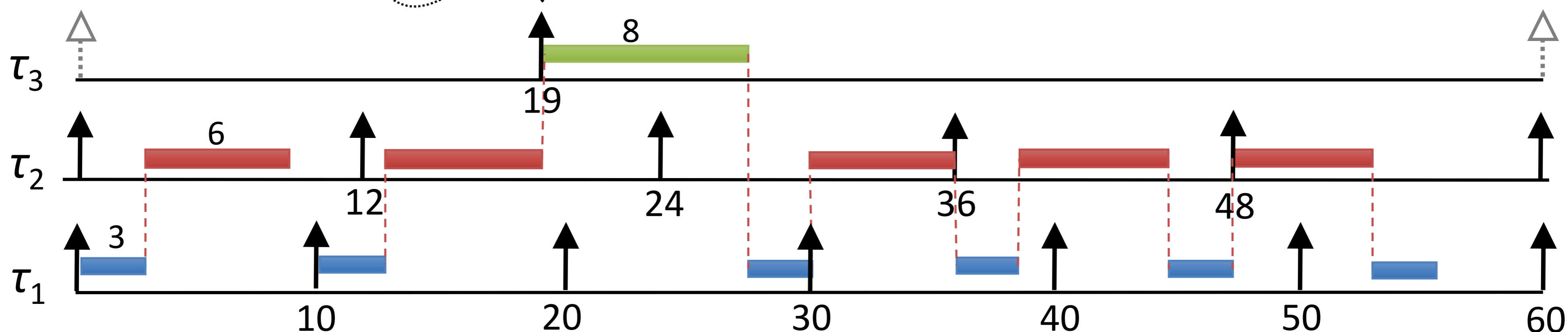
Move  $\tau_3$  "out of the way" by *introducing* (or *adjusting*) a *release offset*.

***FIFO schedule becomes identical to CW-EDF schedule!***

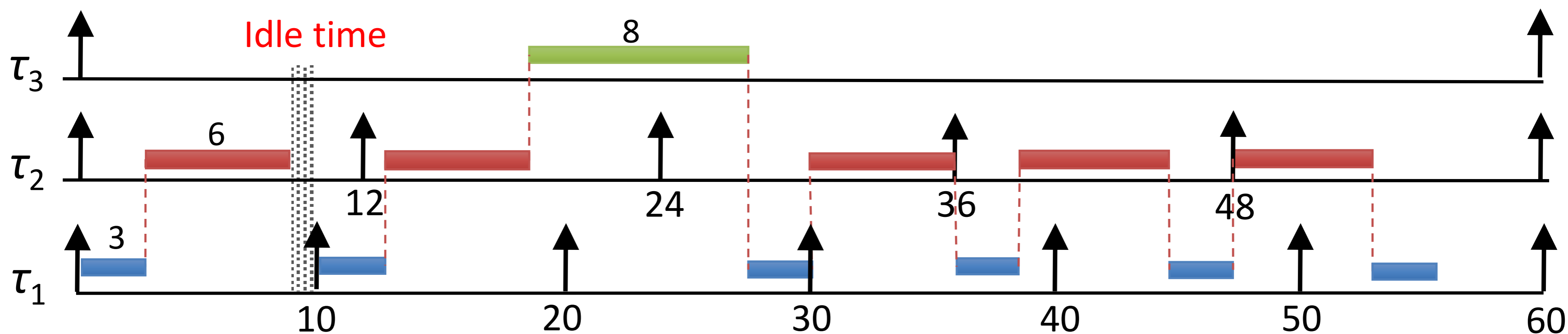
# INTUITION: FIFO + "JUST THE RIGHT" OFFSETS

**FIFO** schedule + offset for  $\tau_3$ :

[Altmeyer, Sundharam, & Navet, 2016]



**CW-EDF** schedule is identical:



**THIS PAPER**

***OFFSET TUNING ALGORITHM***

# PROBLEM STATEMENT

Given a set of  $n$  *periodic non-preemptive tasks*, find, for each job of each task, a **release offset** such that

(A) the resulting **FIFO schedule is feasible**, and

(B) the **number of offsets** per task is **minimized**.

## Challenges

- space of possible offsets is large and unstructured
- even ignoring (B), solving "just" (A) is *very* difficult

## Altmeyer et al.

- randomize offsets + test
- not systematic
- scalability limitations

[S. Altmeyer, S. Sundharam, and N. Navet, "The case for FIFO real-time scheduling," University of Luxembourg, Tech. Rep., 2016]

## KEY INSIGHT

Given a set of  $n$  *periodic non-preemptive tasks*, find, for each job of each task, a **release offset** such that

(A) the resulting **FIFO schedule is feasible**, and

(B) the **number of offsets** per task is **minimized**.

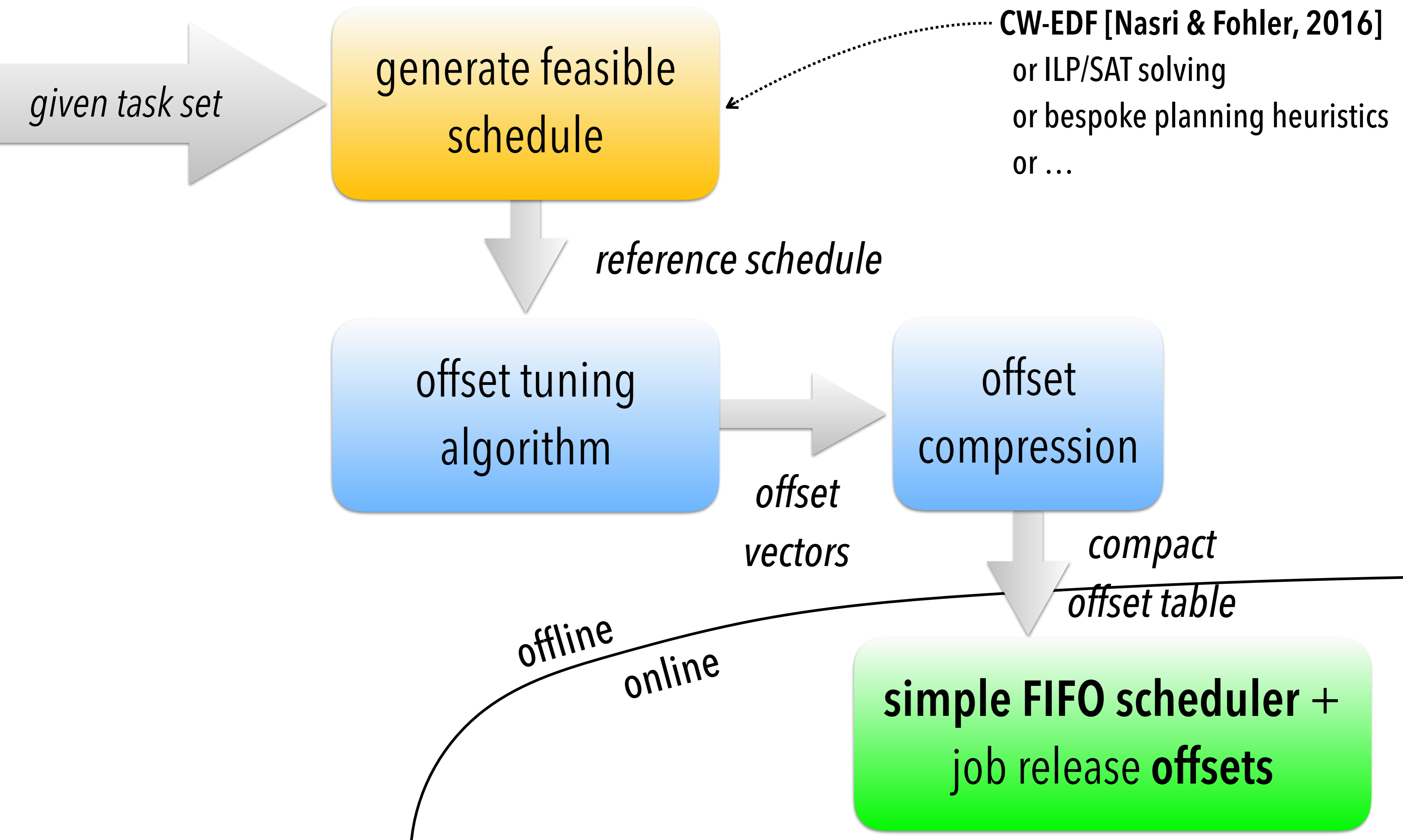
*Solving (A) is very difficult... so we don't!*

## OFFSET TUNING

**Infer offsets** from a given **feasible reference schedule**, while greedily working towards (B).



# OFFSET TUNING – OVERVIEW



# SCHEDULE EQUIVALENCY

A schedule  $S_1$  is **equivalent to**  $S_2$  if

- (i) they schedule the *same jobs*,
- (ii) in the *same order*, and
- (iii) *jobs start no later* in  $S_1$  than in  $S_2$ .

## Non-preemptive execution

→ jobs also complete no later in  $S_1$  than in  $S_2$

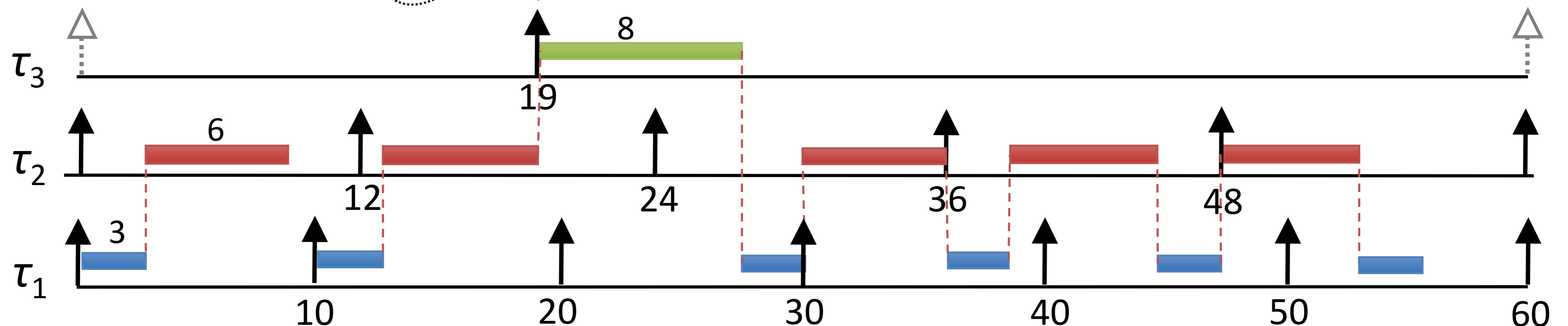
## Offset Tuning

→ ensures FIFO schedule is equivalent to reference schedule

# POI: POTENTIAL OFFSETS INTERVAL

**POI of a job:** range of release offsets that *guarantee schedule equivalency*.

**FIFO** schedule + offset for  $\tau_3$ :

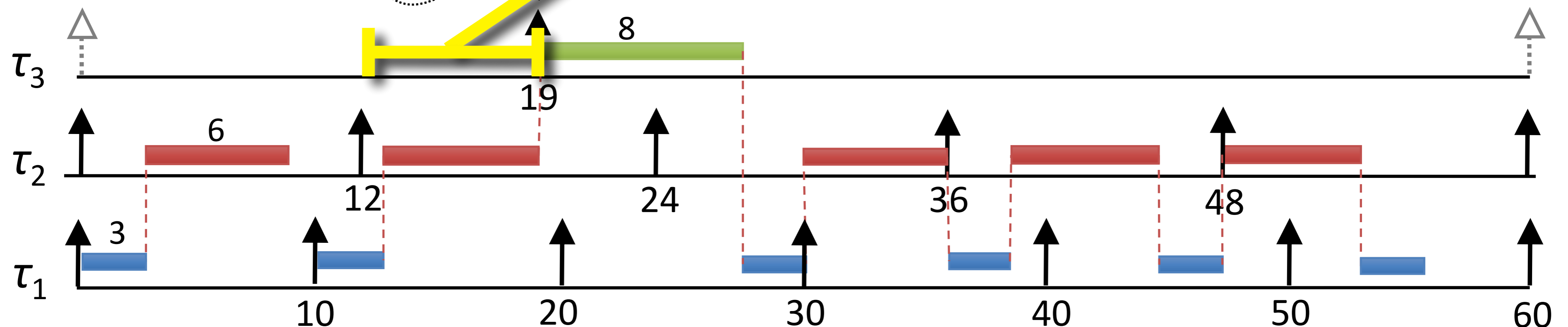


# POI: POTENTIAL OFFSETS INTERVAL

**POI of a job:** range of release offsets that *guarantee schedule equivalency*.

**POI:** any release time of  $\tau_3$  in **(12, 19]** will yield an **equivalent schedule**.

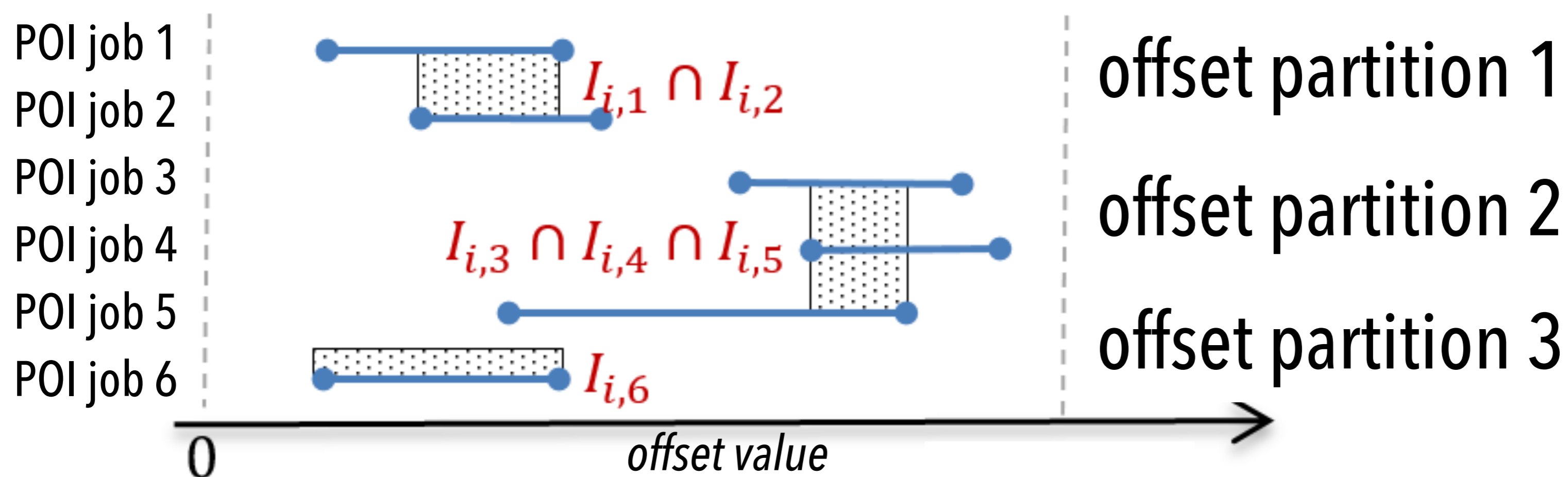
**FIFO** schedule + *offset* for  $\tau_3$ :



# OFFSET PARTITION

Consecutive jobs of a task form an **offset partition** if they have *mutually intersecting POIs*.

→ can be assigned a single offset



→ offset partitioning not necessarily unique

# OFFSET TUNING ALGORITHM (*SIMPLIFIED*)

for each task  $\tau_i$  in deadline-monotonic order:

**greedily** create *offset partitions* for  $\tau_i$

**assuming** *jobs of larger-deadline tasks are released as in reference schedule*

Need to start somewhere...

**shorter relative deadline = fewer options**

## ALGORITHM (**SIMPLIFIED**)

for each task  $\tau_i$  in deadline-monotonic order:

**greedily** create *offset partitions* for  $\tau_i$

**assuming** *jobs of larger-deadline tasks are released as in reference schedule*

Release times of not-yet-processed jobs still unknown → **speculate**.

*Mis-speculation increases the number of offset partitions,  
but **does not** cause the algorithm to fail.*

# PROPERTIES OF OFFSET TUNING

## REFERENCE SCHEDULE EQUIVALENCY

*In the resulting FIFO schedule, no job completes later than in the original reference schedule.*

## PER-TASK MINIMAL OFFSET PARTITIONS

*The greedy offset partitioning strategy yields a minimal number of offset partitions (for a given task).*

## NON-MINIMAL OFFSET PARTITIONS FOR ENTIRE TASK SET

*Deadline-monotonic processing order does not guarantee overall minimal number of offset partitions (but **works well empirically**).*



# SINGLE-OFFSET HEURISTICS

**What if we want just a *single offset* per task?**

- no extra memory required
- compatibility with existing systems

**FST: First-Start-Time Heuristic**

- pick start time of first job in reference schedule

**FOP: First-Offset-Partition Heuristic**

- pick offset from first offset partition of the task

# **EVALUATION**

# EVALUATION QUESTIONS

**Q1: Does FIFO + Offset Tuning still have low runtime overheads?**

**Q2: Does FIFO + Offset Tuning (FIFO-OT) significantly improve schedulability relative to EDF/RM?**

**Q3: How many offsets are assigned?**

**Q4: How much memory is needed?**

# PROTOTYPE PLATFORM

## Arduino Mega 2560

ATMega2560 microcontroller

16 MHz CPU

256 KiB Flash

8 KiB SRAM (no cache)



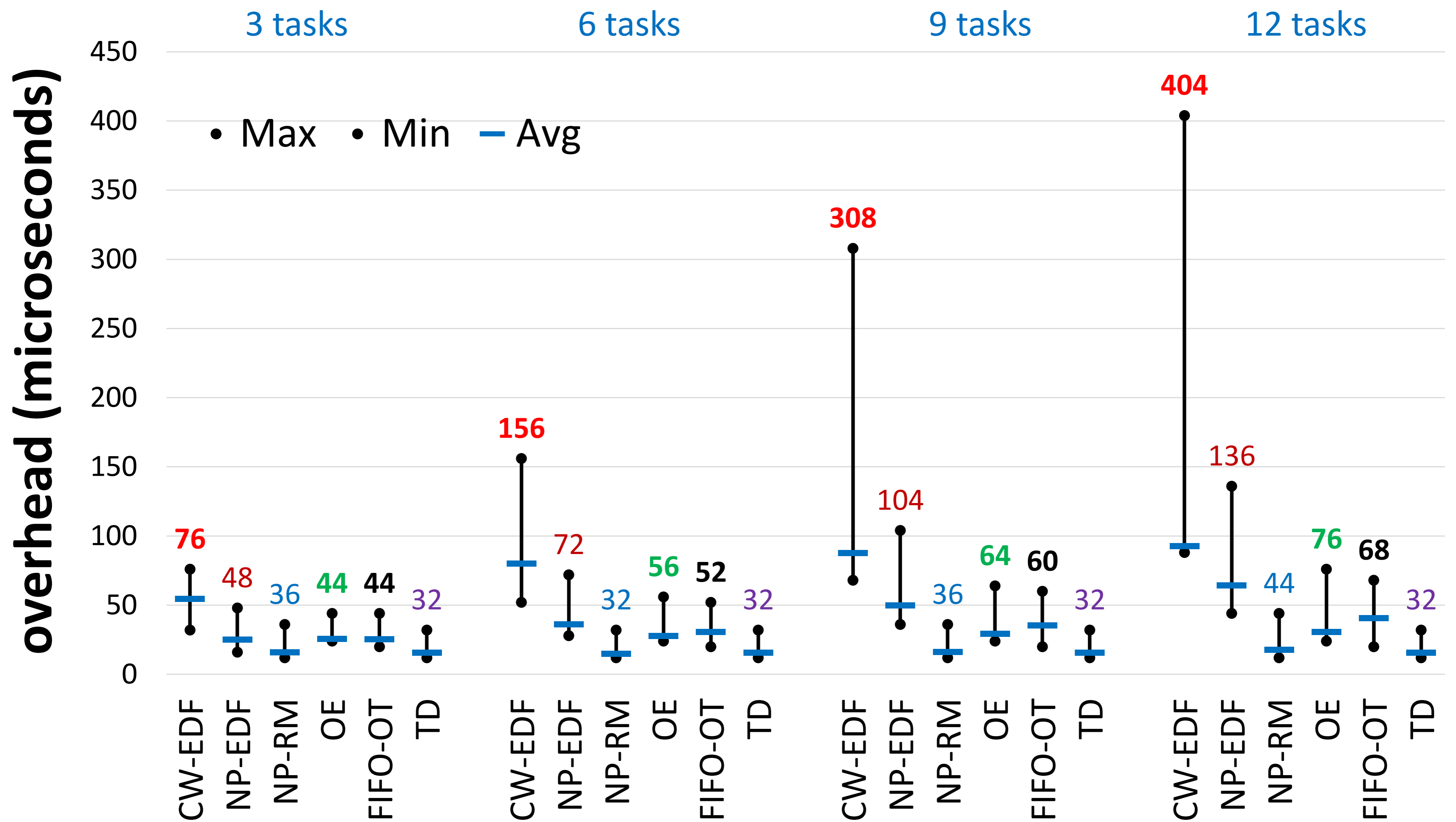
gcc: -Os

<http://people.mpi-sws.org/~bbb/papers/details/rtas18>

# EVALUATED SCHEDULERS

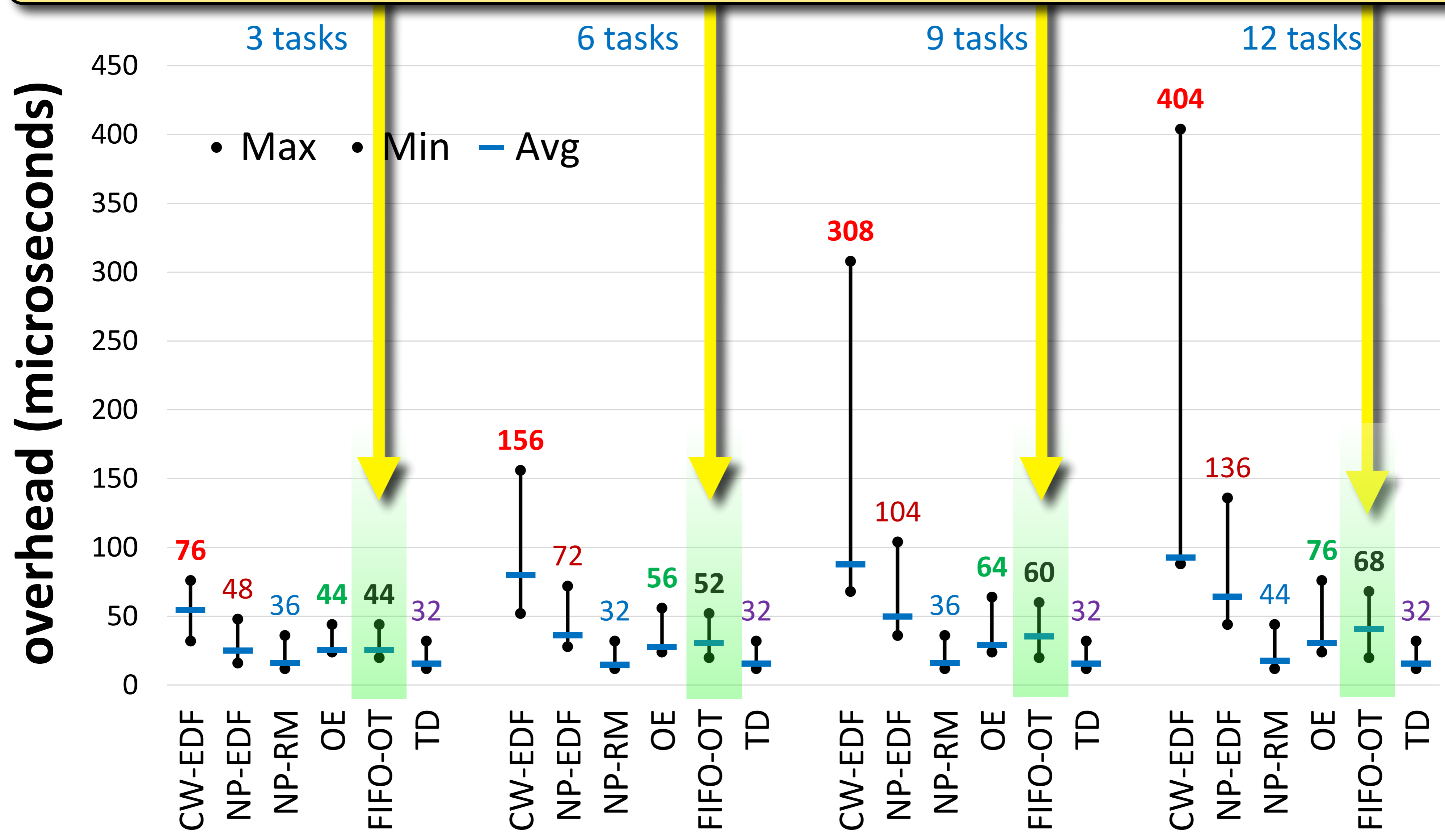
<b>NP-RM</b>	plain non-preemptive rate-monotonic scheduling
<b>NP-EDF</b>	plain non-preemptive EDF
<b>CW-EDF</b>	Critical Window EDF [ <i>Nasri &amp; Fohler, 2016</i> ]
<b>TD</b>	Table-driven (a.k.a. static or time-triggered) scheduling
<b>OE</b>	Offline Equivalence [ <i>Nasri &amp; Brandenburg, 2017</i> ]
<b>FIFO-OT</b>	FIFO + Offset Tuning [ <i>this paper</i> ]

# Q1: RUNTIME OVERHEADS



# LOW RUNTIME OVERHEADS

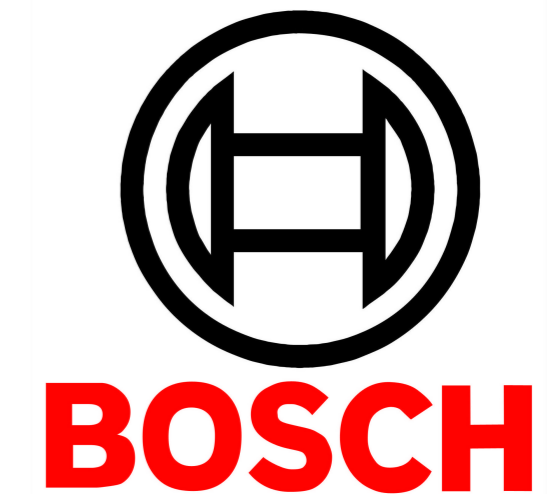
**FIFO-OT is much cheaper than CW-EDF and roughly similar to NP-RM and OE.**



# WORKLOADS

*based on*

Kramer, Ziegenbein, and Hamann, "Real world automotive benchmark for free," WATERS 2015



## **Periods**

→ non-uniformly in {1, 2, 5, 10, 20, 50, 100, 200, 1000} milliseconds

## **Runnable BCETs and WCETs**

→ randomly generated based on statistics provided by Kramer et al.

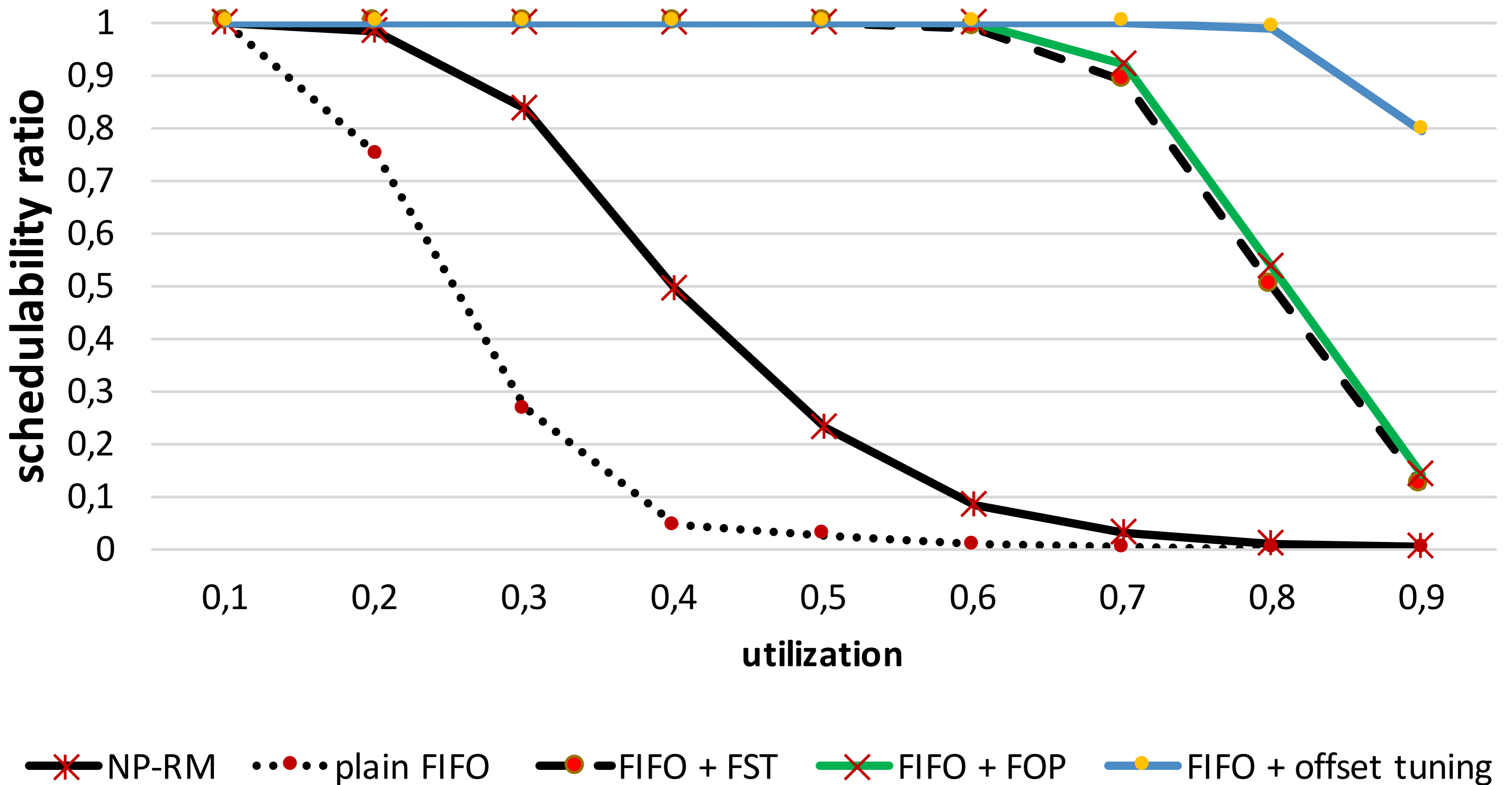
## **Runnable Packing**

→ Runnables aggregated into tasks until random utilization threshold reached

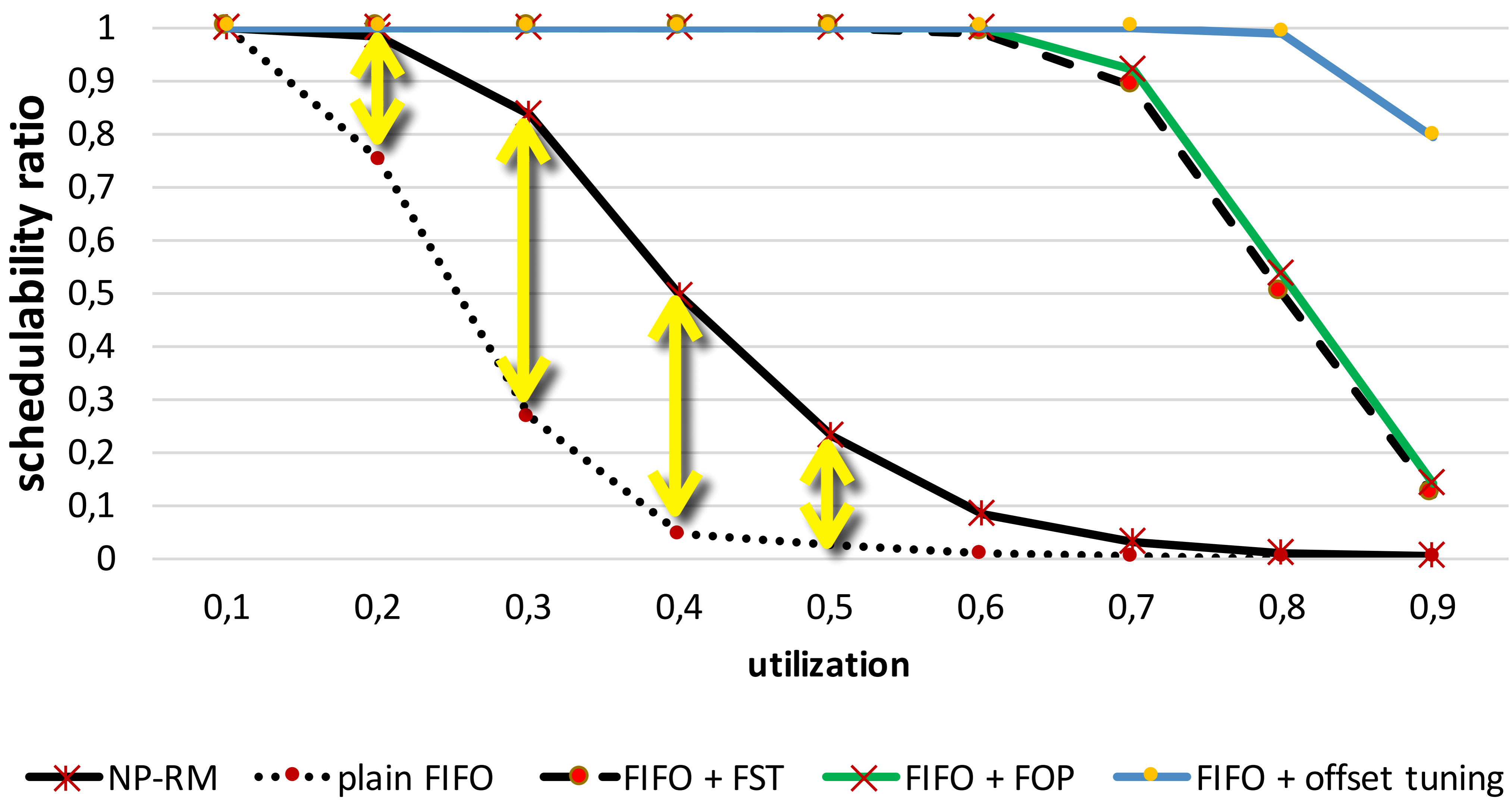
→ *utilization threshold ensures feasibility under non-preemptive scheduling*



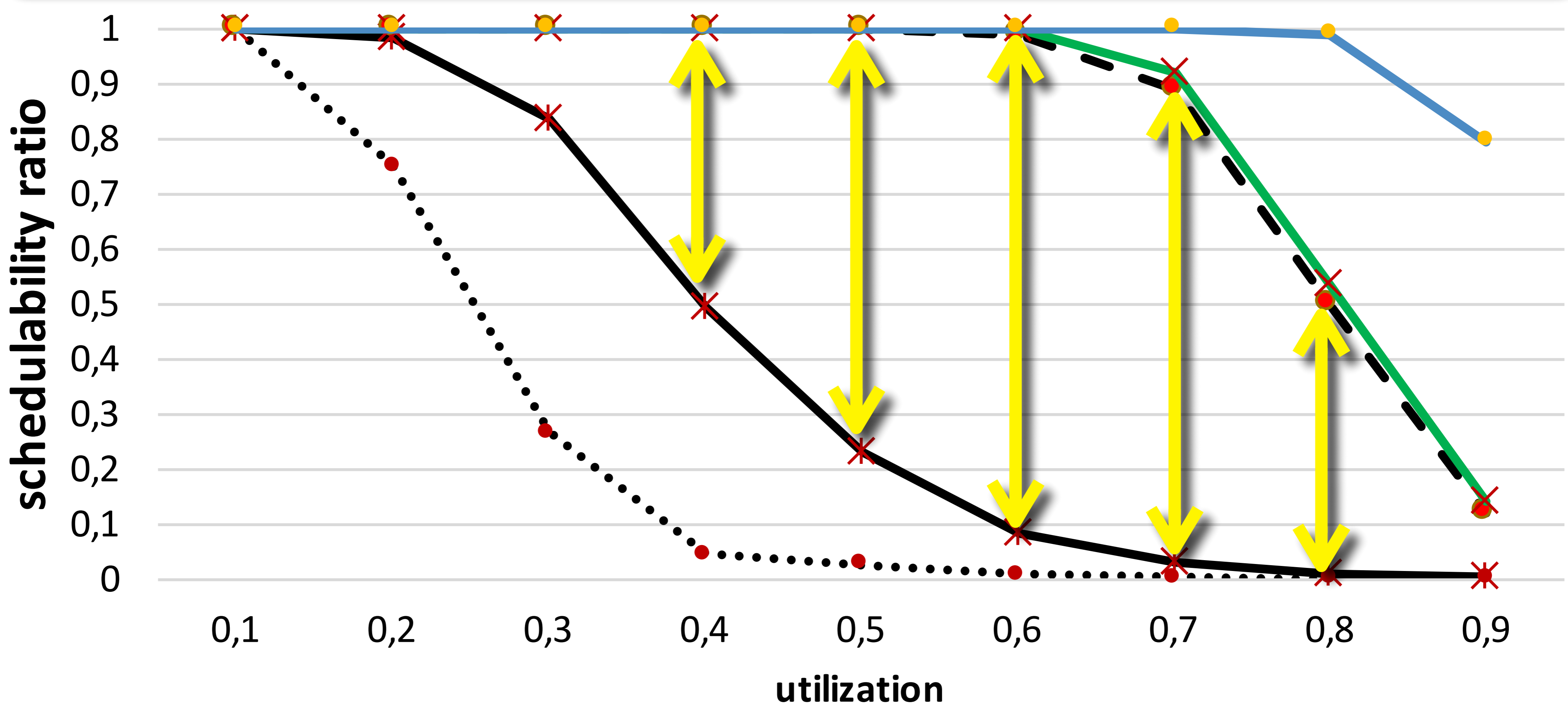
# Q2: SCHEDULABILITY GAINS



As expected, **plain FIFO** exhibits very low schedulability.

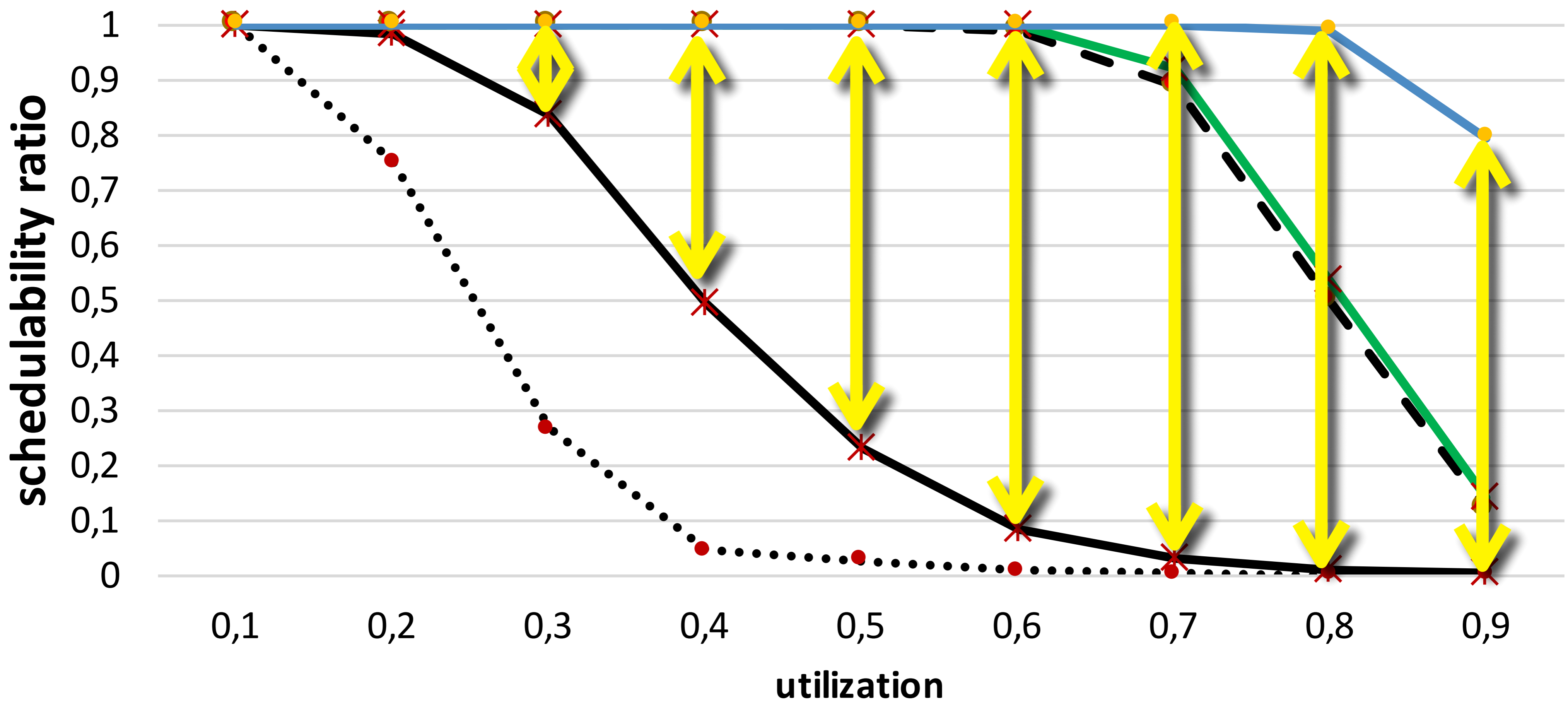


Assigning even a *single offset* per task can substantially increase schedulability!



—x— NP-RM   
 ···●··· plain FIFO   
 —●— FIFO + FST   
 —x— FIFO + FOP   
 —●— FIFO + offset tuning

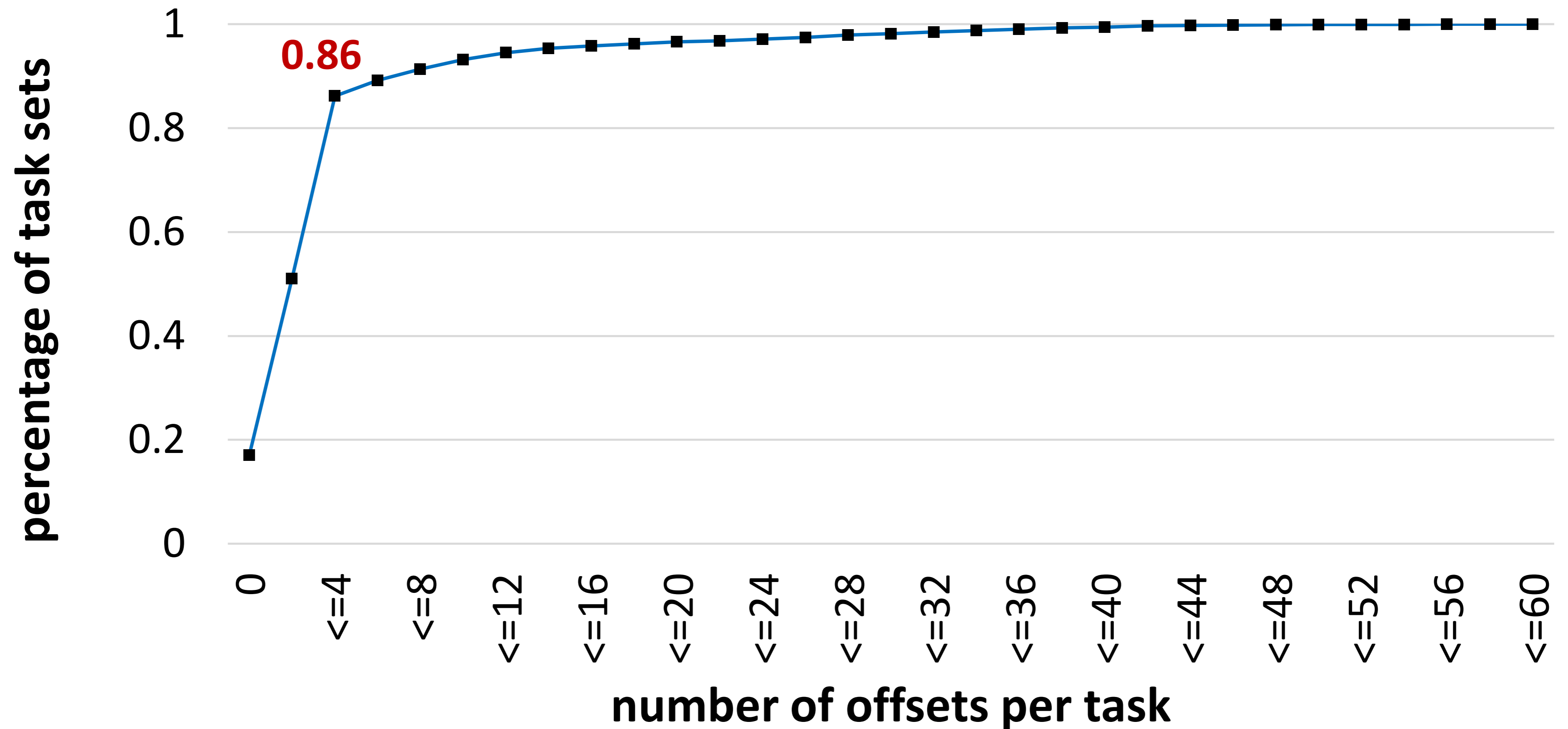
## Q2: SCHEDULABILITY GAINS



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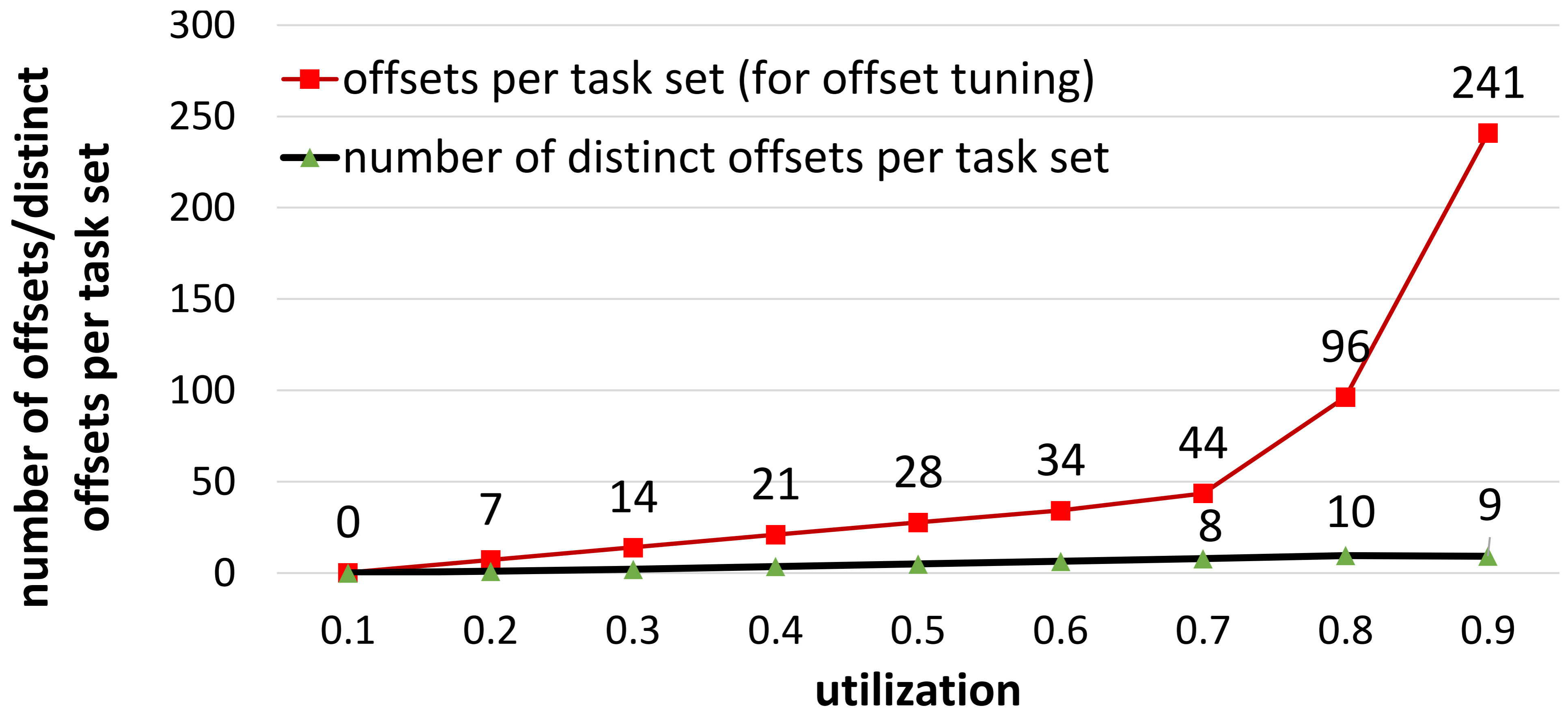
**FIFO-OT achieves *much higher schedulability*, thanks to CW-EDF reference schedule.**

## Q3: NUMBERS OF OFFSETS PER TASK



→ *Most tasks require only few offset partitions.*

# NUMBERS OF **UNIQUE** OFFSETS PER TASK SET

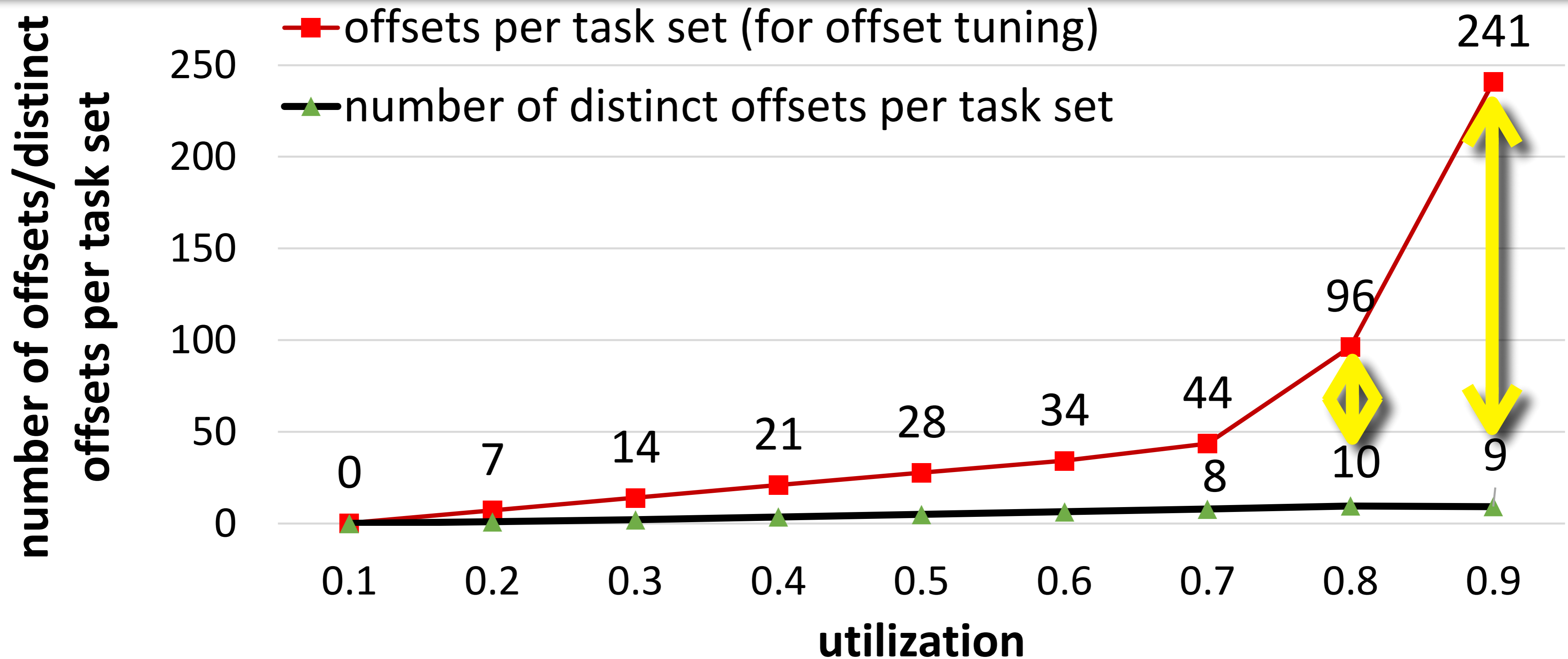


Across the hyper-period, **offsets values repeat cyclicly**.

→ *Opportunity to store offsets efficiently (compression).*

# NUMBERS OF **UNIQUE** OFFSETS PER TASK SET

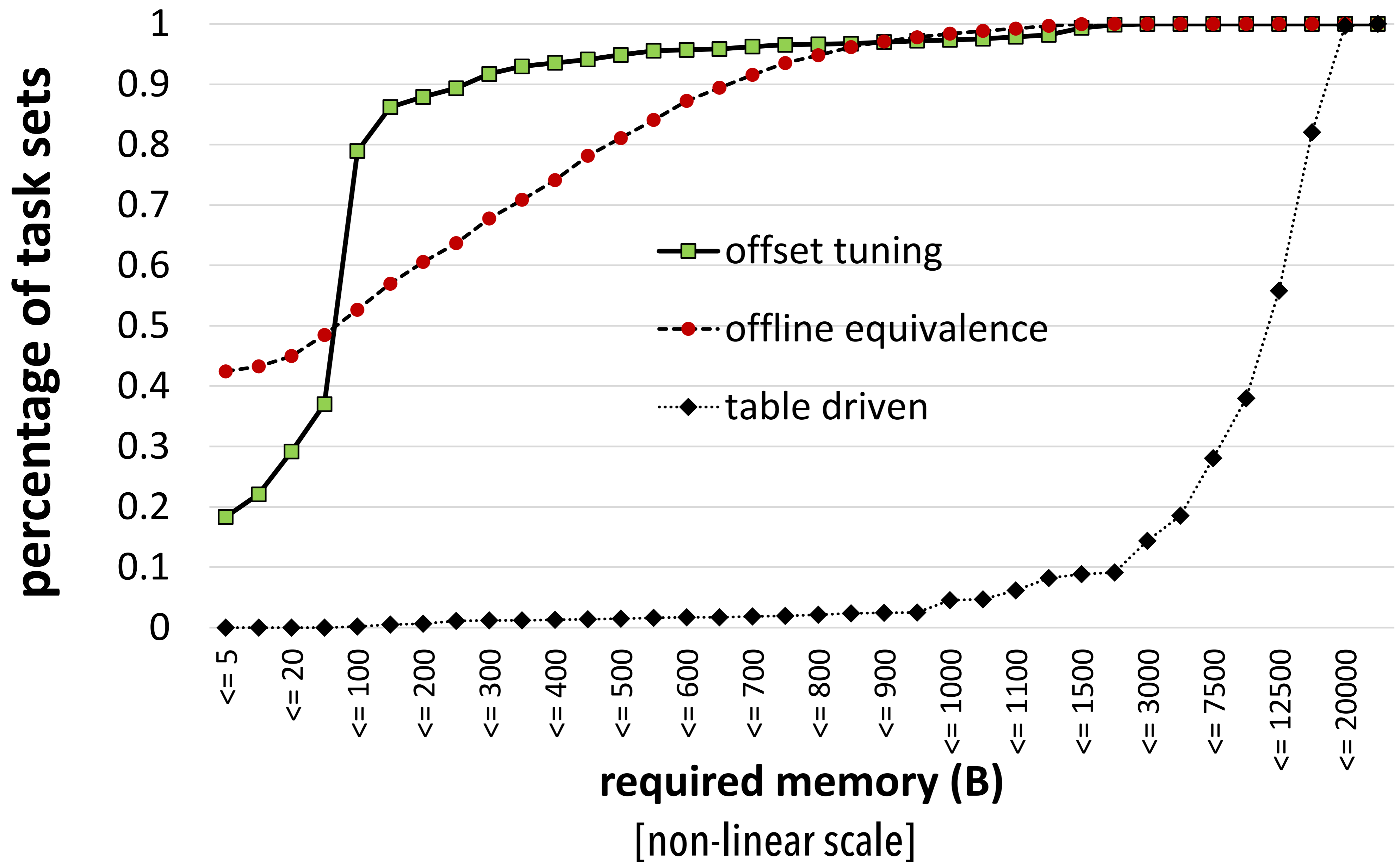
Up to **25× reduction** in the number of offset values that must be stored.



Across the hyper-period, **offsets values repeat cyclicly**.

→ Opportunity to store offsets efficiently (compression).

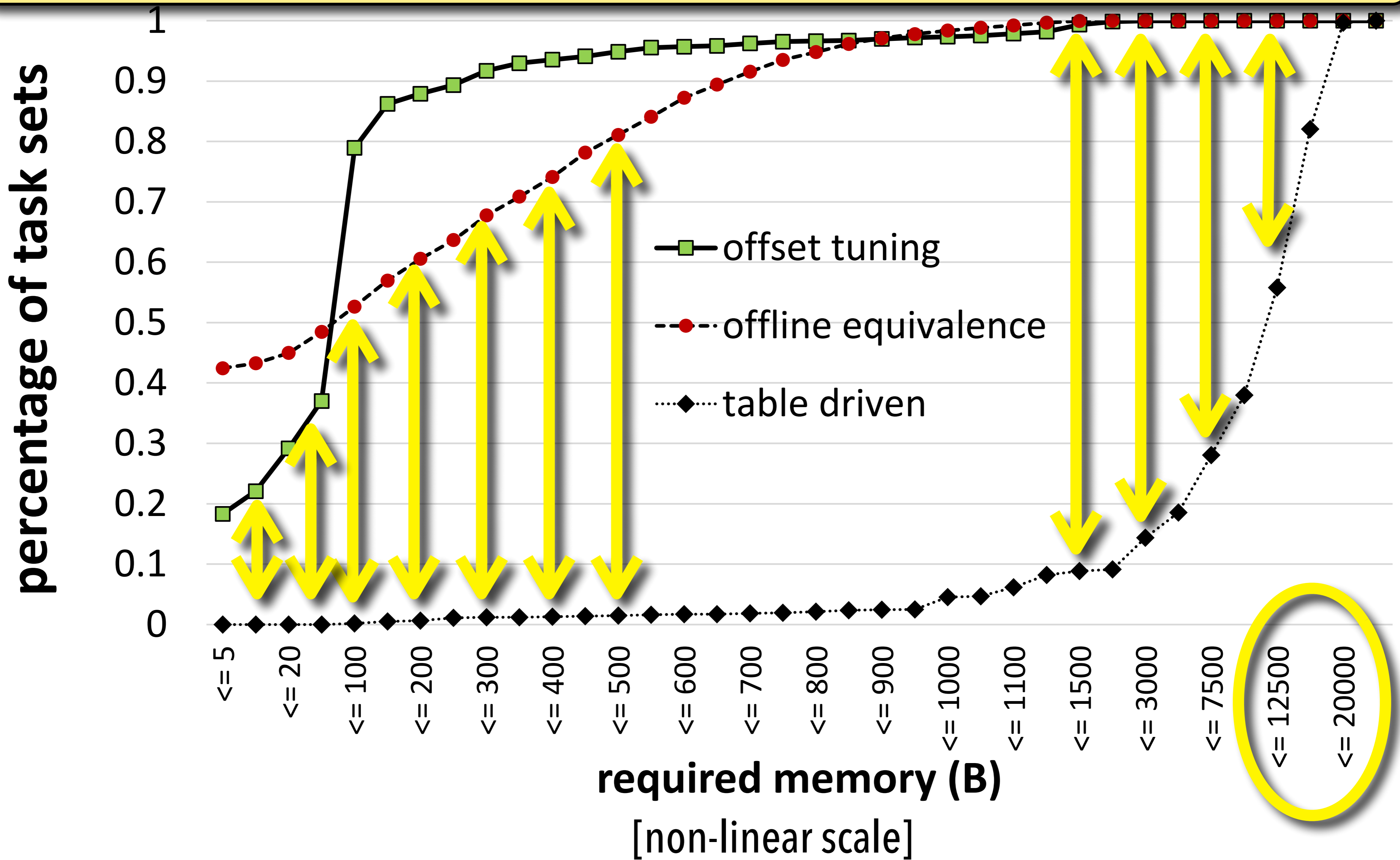
# MEMORY USAGE





# Both *OE* and *FIFO-OT* require much less memory than table-driven scheduling.

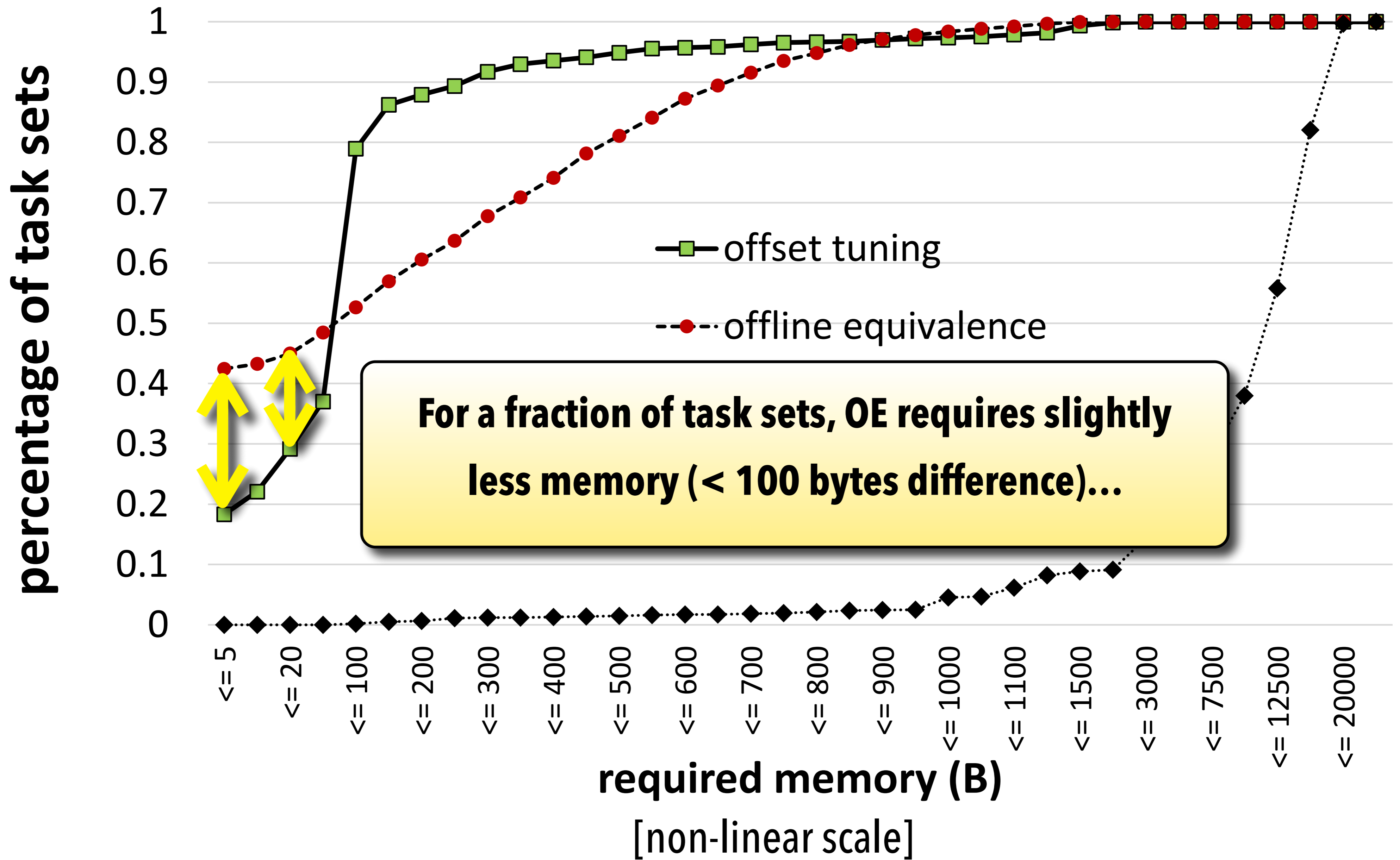
*dozens to hundreds of bytes vs. 10KiB-20KiB*



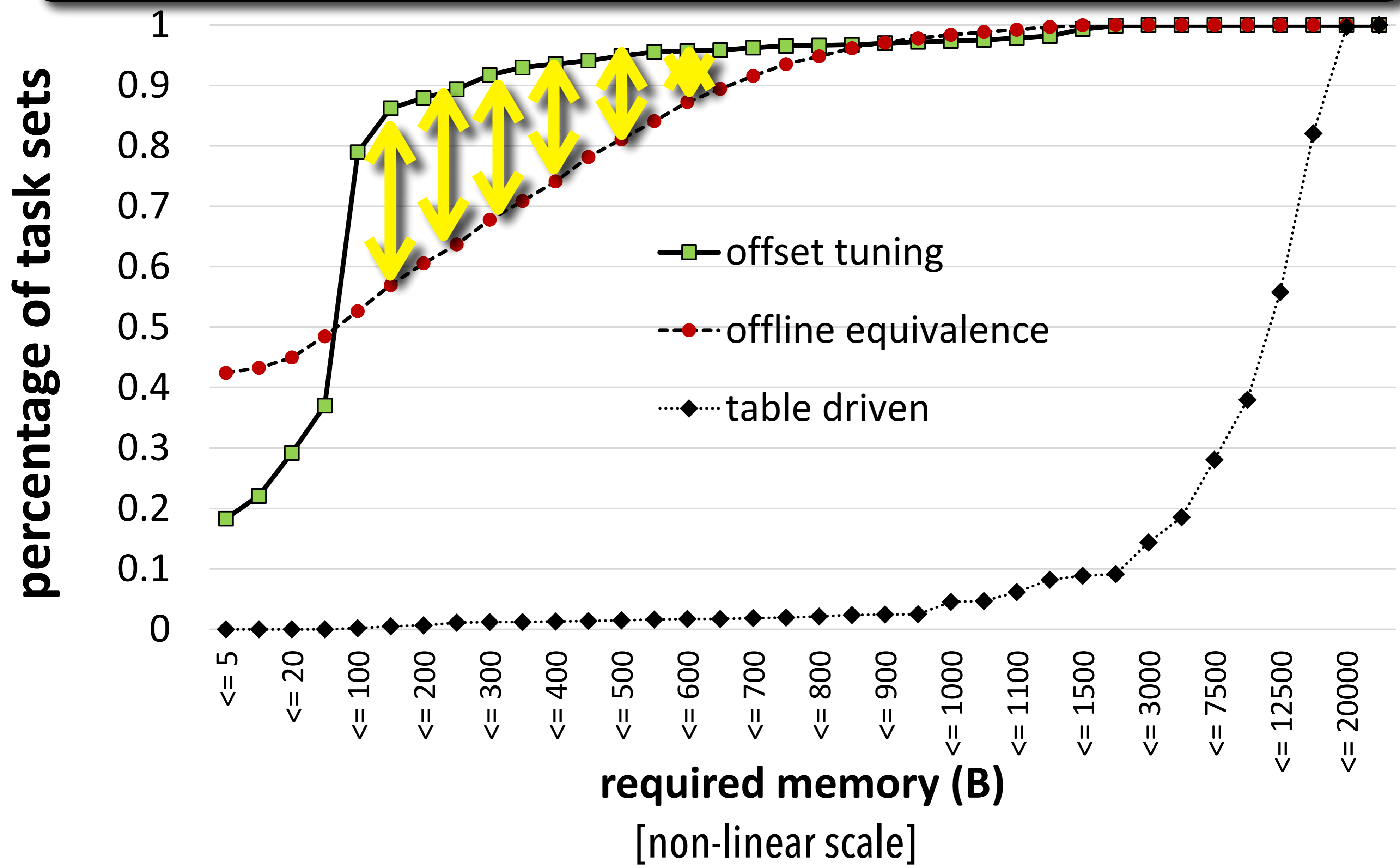
<= 12500

<= 20000

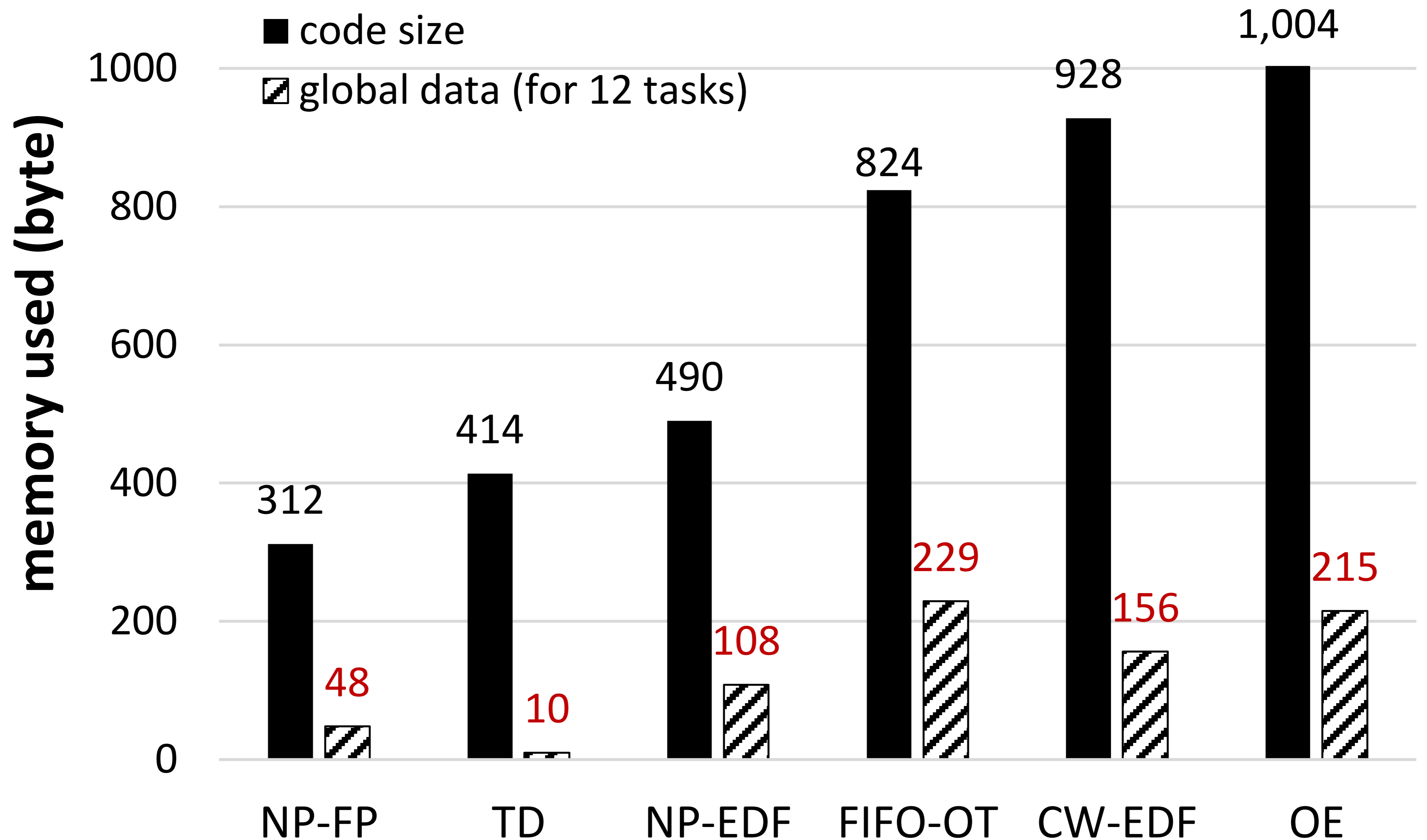
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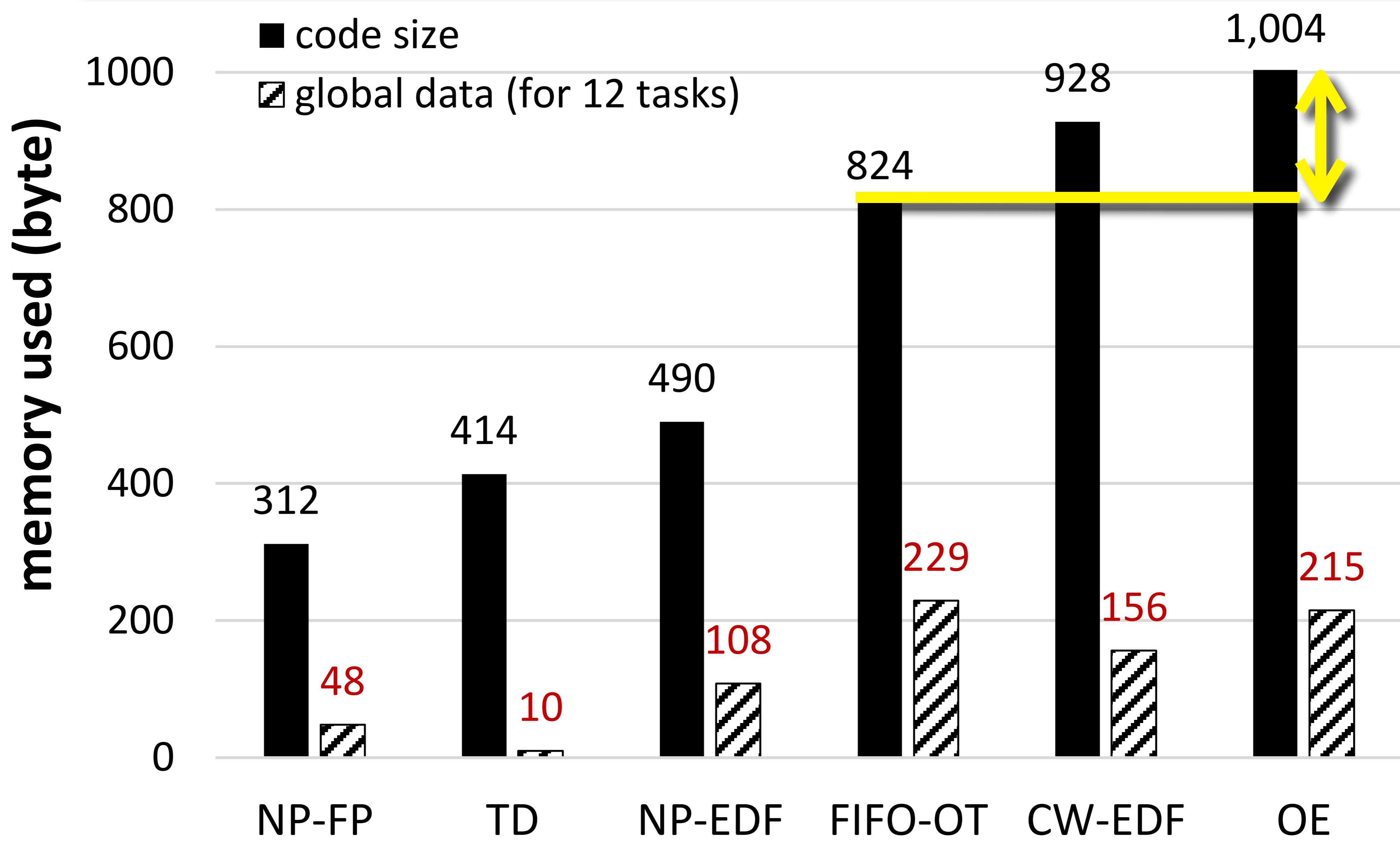
...but FIFO-OT can support over *90% of task sets with  $\leq 250$  bytes of offset data.*



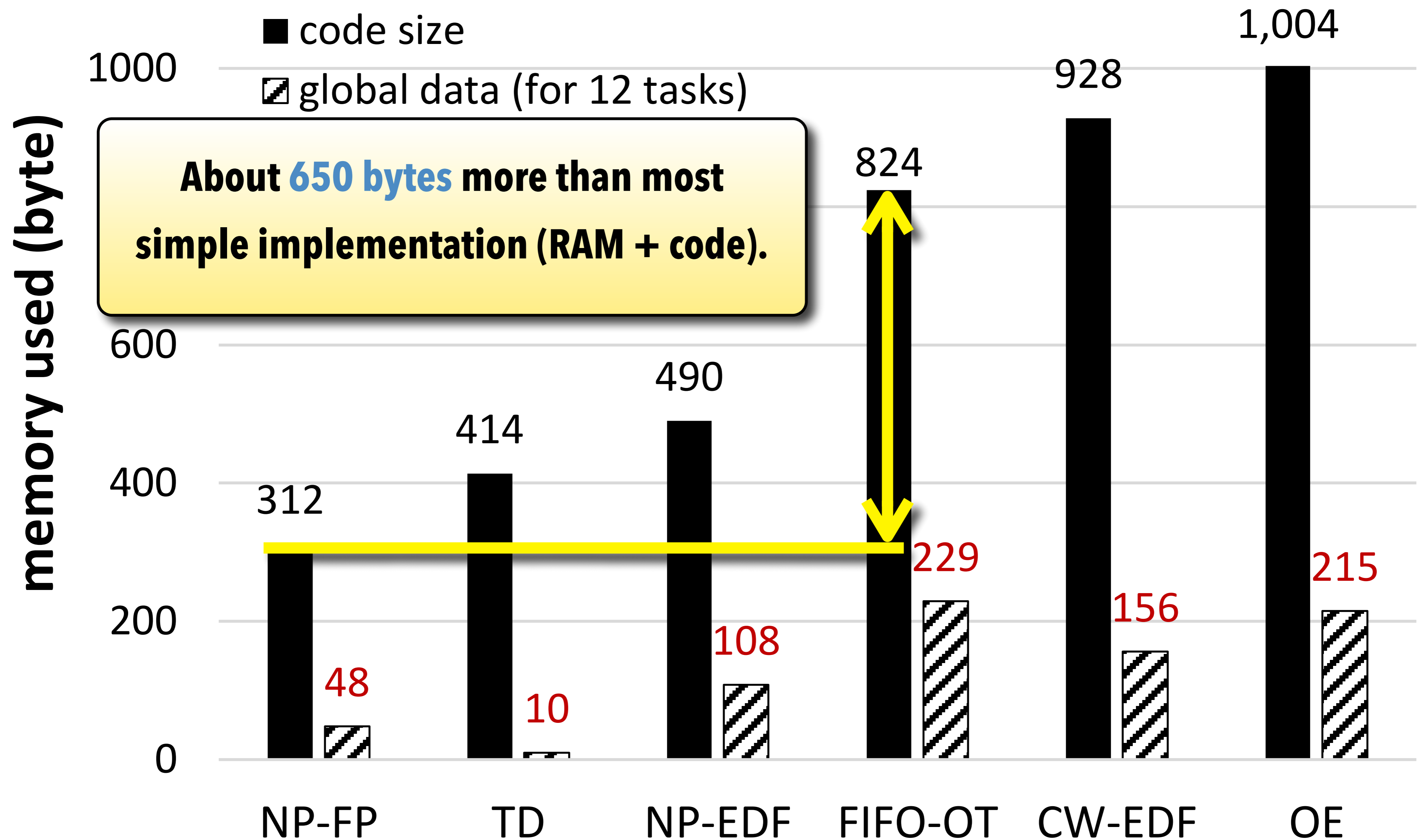
# IMPLEMENTATION FOOTPRINT



About 150 bytes *smaller footprint* than OE (RAM + code).



# IMPLEMENTATION FOOTPRINT



# **CONCLUSION**

# FIFO SCHEDULING

First-In-First-Out (**FIFO**) scheduling

- extremely simple
- very low overheads

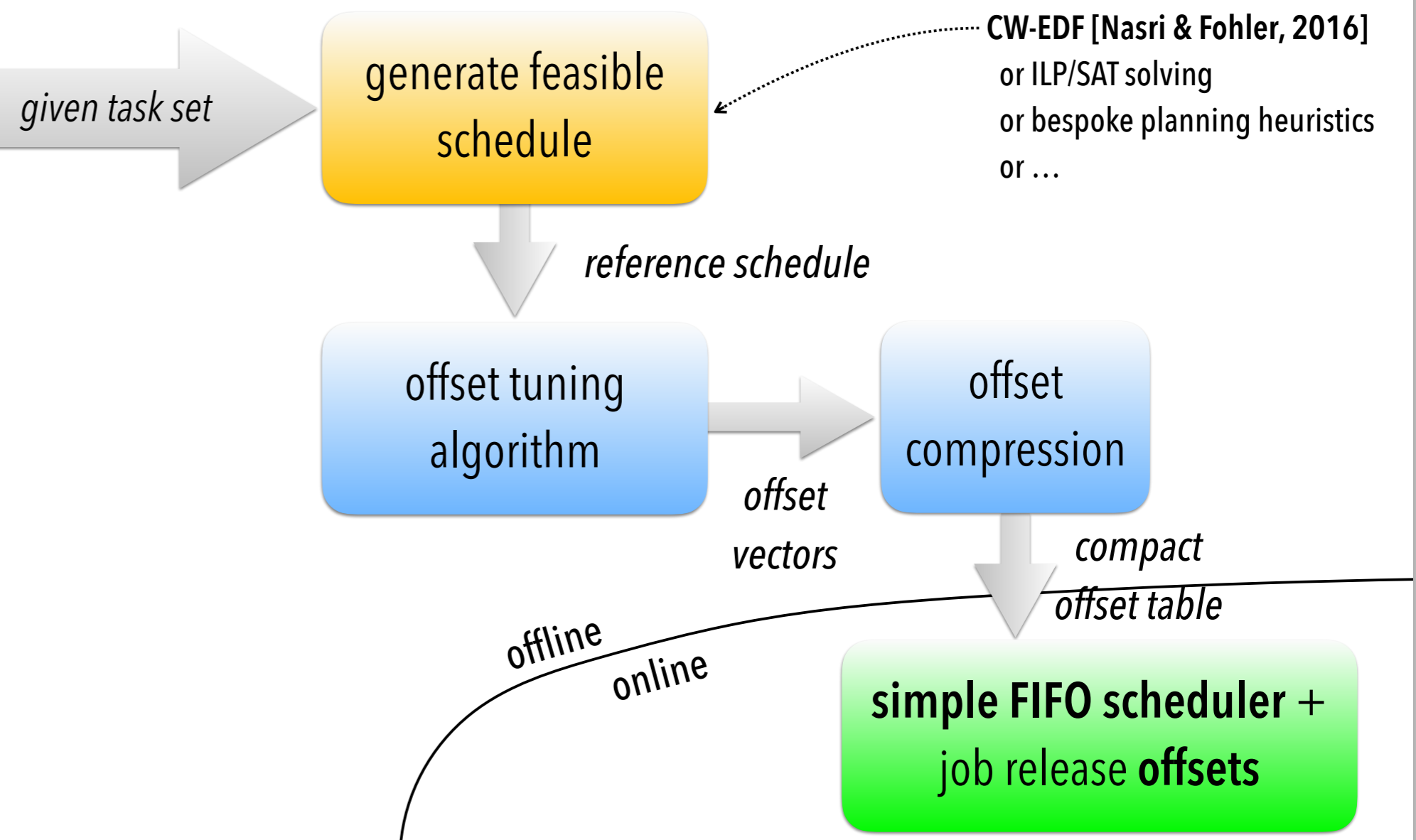
ideal for:   
 IoT-class devices   
 deeply embedded systems   
 hardware implementations

very ~~low~~ **schedulability** } meeting **deadlines?**   
 **HIGH!**

## THIS PAPER

**FIFO can actually achieve excellent schedulability!**   
 [periodic non-preemptive tasks on a uniprocessor]

# OFFSET TUNING – OVERVIEW



# PROPERTIES OF OFFSET TUNING

## REFERENCE SCHEDULE EQUIVALENCY

In the resulting FIFO schedule, no job completes later than in the original reference schedule.

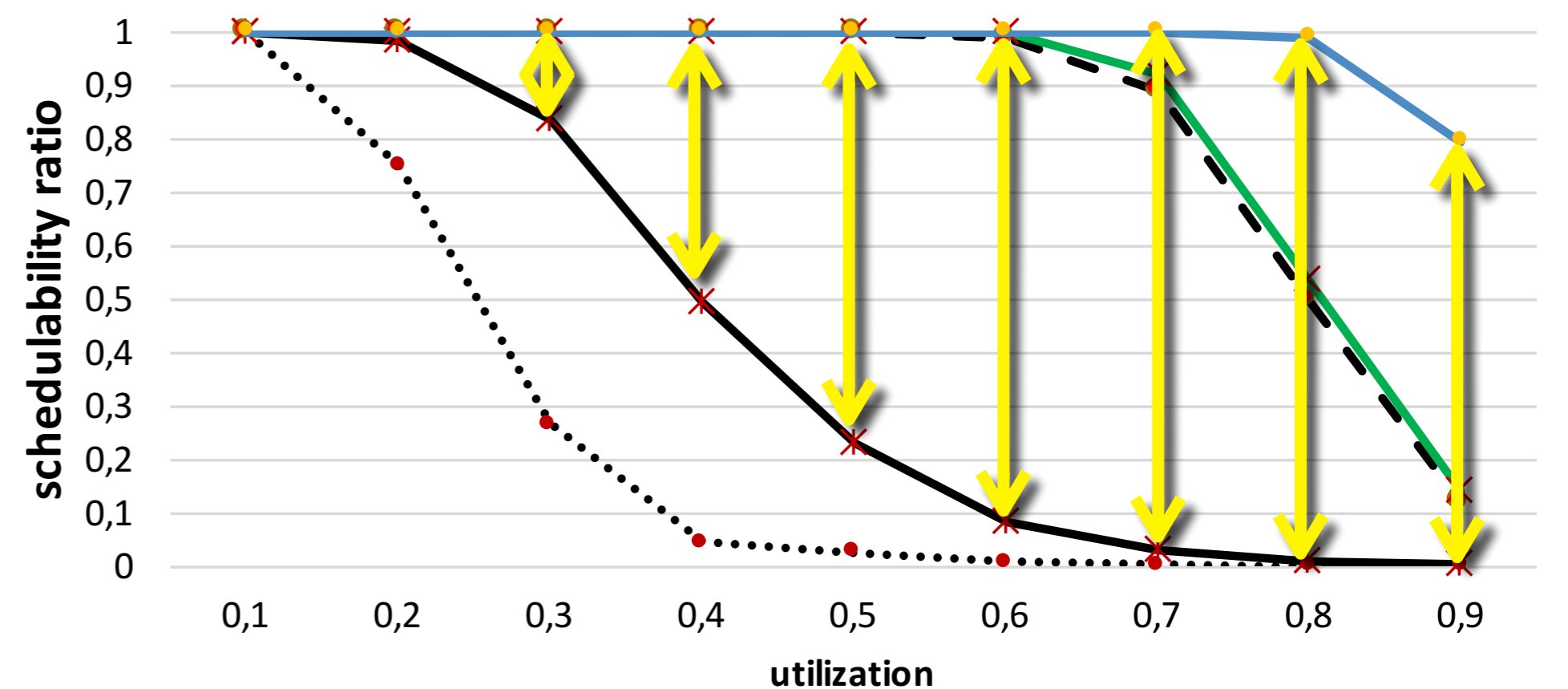
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The greedy offset partitioning strategy yields a minimal number of offset partitions (for a given task).

## NON-MINIMAL OFFSET PARTITIONS FOR ENTIRE TASK SET

Deadline-monotonic processing order does not guarantee overall minimal number of offset partitions (but **works well empirically**).

# Q2: SCHEDULABILITY GAINS



—x— NP-RM    ··· plain FIFO    —•— FIFO + FST    —x— FIFO + FOP    —◇— FIFO + offset tuning

**FIFO-OT achieves much higher schedulability, thanks to CW-EDF reference schedule.**

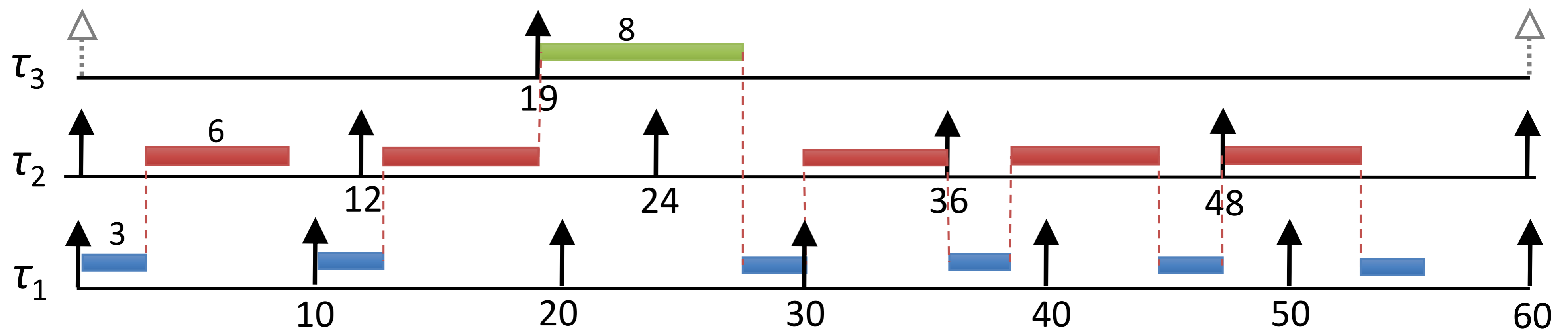


# **APPENDIX**

# CAN OFFSET TUNING BE APPLIED TO EDF OR FIXED-PRIORITY SCHEDULING?

→ *yes in principle, but no equivalence guarantee*

**FIFO** schedule + offset for  $\tau_3$ :



**RM** schedule + offset for  $\tau_3$ :

